

## Outline

- Chapter 4: Processes
- Chapter 6: CPU Scheduling
- Chapter 5: Threads
- Introduction to Threads: Birrell
- Continuations: Draves et al.



## Processes

- Process is a program in execution
  - Program code
  - Data section
  - Execution context: Program counter, registers, stack
- Process has thread(s) of control
- Many processes “run” concurrently: Process scheduling
  - Fair allocation of CPU and I/O bound processes
    - Context switch



## Process creation

- Creating new processes is expensive
  - Resource allocation issue
- Fork mechanism: UNIX, Windows NT
  - Duplicate the parent process
  - Shares file descriptors, memory is copied
  - Exec to create different process
  - Various optimizations to avoid copying the entire parent context (Copy on write (COW), etc..)
- Exec mechanism: VMS, Windows NT
  - New process is specifically loaded



## Interprocess communication

- Processes need to communicate with each other
  - Naming
  - Message-passing
    - Direct (to process) or indirect (port, mailbox)
    - Symmetric or asymmetric (blocking, nonblocking)
    - Automatic or explicit buffering (capacity)
    - Send by copy or reference
    - Fixed size or variable size messages
  - Shared memory/mutexes
  - Remote Procedure Call (RPC/RMI)



## CPU scheduling

- Interleave processes so as to maximize utilization of CPU and I/O resources
- Scheduler should be fast as time spent in scheduler is wasted time
  - Switching context (h/w assists – register windows [sparc])
  - Switching to user mode
  - Jumping to proper location
- Preemptive scheduling:
  - Process could be in the middle of an operation
  - Especially bad for kernel structures
- Non-preemptive (cooperative) scheduling:
  - Starvation



## Classical algorithms

- First come, first served (FCFS)
  - Convoy effect (all processes get bunched up behind long process)
- Shortest job first (shortest next CPU burst)
  - “Optimal”
  - Need oracle to predict next duration: prediction based on history
- Priority scheduling
  - Priority inversion
  - Starvation
- Round robin scheduling
- Multilevel queue scheduling
- Multilevel feedback queue scheduling
- Multiprocessor scheduling
  - Gang scheduling, Non Uniform Memory Access, Processor affinity
- Real-Time scheduling
  - Hard and soft real time systems



## Threads

- Applications require concurrency. Threads provide a neat abstraction to specify concurrency
- E.g. word processor application
  - Needs to accept user input, display it on screen, spell check and grammar check
  - Implicit: Write code that reads user input, displays/formats it on screen, calls spell checked etc. while making sure that interactive response does not suffer. May or may not leverage multiple processors
  - Threads: Use threads to perform each task and communicate using queues and shared data structures
  - Processes: expensive to create and do not share data structures and so explicitly passed



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## Threads - Benefits

- Responsiveness
  - If one “task” takes too long, other “tasks” can still proceed
- Resource sharing:
  - Grammar checker can check the buffer as it is being typed
- Economy:
  - Process creation is expensive (spell checker)
- Utilization of multiprocessor architectures:
  - If we had four processors (say), the word processor can fully leverage them
- Pitfalls:
  - Shared data should be protected or results are undefined
    - Race conditions, dead locks, starvation (more Thurs)



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## Thread types

- Continuum: Cost to create and ease of management
- User level threads (e.g. pthreads)
  - Implemented as a library
  - Fast to create
  - Cannot have blocking system calls
  - Scheduling conflicts between kernel and threads. User level threads cannot do anything is kernel preempts the process
- Kernel level threads
  - Slower to create and manage
  - Blocking system calls are no problem
  - Most OS's support these threads



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## Threading models

- One to One model
  - Map each user thread to one kernel thread
- Many to one model
  - Map many user threads to a single kernel thread
  - Cannot exploit multiprocessors
- Many to many
  - Map  $m$  user threads to  $n$  kernel threads



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## Threading Issues:

- Cancellation:
  - Asynchronous or deferred cancellation
- Signal handling:
  - Relevant thread
  - Every thread
  - Certain threads
  - Specific thread
- Pooled threads (web server)
- Thread specific data



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## Threads – Andrew Birrell

- Seminal paper on threads programming
  - Old but most techniques/experiences are still valid
- Birrell
  - Xerox PARC □ Dec SRC □ Microsoft Research
  - Invented Remote Procedure Calls (RPC)
  - Personal Juke box (hard disk based mp3 – Apple iPod?)
  - Worked on Cedar, Distributed FS etc for ~25 years (1977-



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## Threads – Andrew Birrell

- Dec SRC and Xerox PARC were the premium systems research labs
- PARC researchers invented:
  - Personal computers - Alto
  - Mouse
  - Windows - Star
  - Bitmapped terminals
  - Icons
  - Ethernet
  - Smalltalk
  - Bravo – first WYSIWYG program
  - Laser printer
  - ...



## Continuations

- Richard Draves – Leads Systems and Networking group at Microsoft Research
  - Mach, Rialto
- Brian Bershad – Univ. of Washington □ CMU □ UW
  - Perennial SOSP contributor....
- Richard F. Rashid – CMU □ Head, MS Research
  - Mach, ....
- Randall Dean – CMU, VP, Mercury Systems?
  - Mach, ?



## Key idea

- Optimizing Mach
  - Application level representation of state while blocked
  - Application code to restore stack
    - Tradeoff stack space for complexity (application code)
- Is this relevant?
  - Current processors have lots of memory, so why bother?
  - Per processor kernel stack
    - reduce cache and TLB misses
- Software engineering concerns?
- Interrupt driven, co-routine style services
  - Much current work in MS Research and other places



## Windows 2000

Image Name	PID	CPU	CPU Time	Mem Usage	Base Pri	Threads
System Idle Process	0	99	73:39:58	16 K	N/A	1
System	4	00	0:01:18	28 K	Normal	38
smss.exe	164	00	0:00:00	32 K	High	6
winlogon.exe	184	00	0:00:03	448 K	High	16
csrss.exe	188	00	0:00:44	216 K	High	10
services.exe	236	00	0:00:58	3,792 K	Above Normal	31
lsass.exe	240	00	0:00:01	1,008 K	Above Normal	13
rfhnd.exe	276	00	0:00:03	36 K	Normal	3
SmTPPrv.exe	324	00	0:00:00	176 K	Normal	3
svchost.exe	416	00	0:00:02	2,292 K	Normal	8
spoolsv.exe	444	00	0:00:00	2,112 K	Normal	14
ALMExec.exe	472	00	0:00:00	232 K	Normal	2
svchost.exe	544	00	0:00:08	4,620 K	Normal	27



## wizard.cse.nd.edu status (using top)

load averages: 0.18, 0.09, 0.09 23:53:20  
 118 processes: 113 sleeping, 3 zombie, 2 on cpu  
 Memory: 512M real, 267M free, 129M swap in use, 1260M swap free

PID	USERNAME	THR	PRI	NICE	SIZE	RES	STATE	TIME	CPU	COMMAND
17213	surendar	1	20	0	1600K	992K	cpu/0	0:08	16.94%	yes
16623	root	1	58	0	4800K	3688K	sleep	0:01	0.19%	sshd
17218	surendar	1	48	0	2112K	1136K	cpu/1	0:00	0.14%	top
16840	www	3	48	0	14M	7096K	sleep	0:00	0.14%	httpd
16626	surendar	1	48	0	3656K	2944K	sleep	0:01	0.08%	tcsh
272	root	21	58	0	4448K	1648K	sleep	1:16	0.06%	syslogd

[...]

- -L option for ps will list the information about lwps



## Discussion

- Constant tension between moving functionality to upper layers; involving the application programmer and performing automatically at the lower layers
- Automatically create/manage threads by compiler/system? (open research question)

